

# Game Theory with Applications to Finance and Marketing, I

## Iterative Deletion of Strictly Dominated Strategies

Consider two players, labeled A and B, in a simultaneous game with their *finite* strategy spaces being respectively

$$S^A = \{a_1, a_2, \dots, a_J\},$$

and

$$S^B = \{b_1, b_2, \dots, b_K\}.$$

Suppose that iterative deletion of strictly dominated strategies<sup>1</sup> in two different sequential orders  $V$  and  $T$  leads to two maximal reductions  $(V_N^A, V_N^B)$  and  $(T_M^A, T_M^B)$  for  $(S^A, S^B)$ , where  $N$  (respectively,  $M$ ) is the number of steps taken to reach the maximal reduction  $(V_N^A, V_N^B)$  (respectively,  $(T_M^A, T_M^B)$ ). We shall prove by contraposition that  $V_N^A = T_M^A$  and  $V_N^B = T_M^B$ . To this end, assume without loss of generality that  $a_j \in V_N^A \cap [T_M^A]^c$ , and we shall demonstrate a contradiction.

Since  $T_M^A$  does not contain  $a_j$ ,  $a_j$  must be strictly dominated by some mixed strategy  $\sigma_A$  and deleted in some step  $m \leq M - 1$  under  $T$ . Can we have  $V_N^B \subset T_m^B$ ? We shall show that the answer is negative.

Suppose instead that  $V_N^B \subset T_m^B$ . If  $\sigma_A$  remains feasible when A is restricted to using pure strategies in  $V_N^A$ , then since  $a_j$  is strictly dominated by  $\sigma_A$ ,  $a_j$  cannot be contained in  $V_N^A$ , which is a contradiction. What if  $\sigma_A$  becomes infeasible when A is restricted to using pure strategies in  $V_N^A$ ? In the latter case, there must exist  $a_i \in T_m^A \cap [V_N^A]^c$  with  $\sigma_A(a_i) > 0$ , but since  $V_N^A$  does not contain  $a_i$ ,  $a_i$  must be strictly dominated by some  $\sigma_A^i$  and deleted at some step  $n_i < N$  under  $V$ . For all  $a_i$  as such, we can replace  $a_i$  by  $\sigma_A^i$  in  $\sigma_A$ , in the sense that we assign probability  $\sigma_A(a_i)$  to  $\sigma_A^i$  rather than  $a_i$ , and call the resulting A's mixed strategy  $\sigma'_A$ . Clearly,  $\sigma_A$  (and hence  $a_j$ ) is

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<sup>1</sup>Let  $S_0^i \equiv S^i$  and  $\Sigma_0^i \equiv \Sigma^i$ . Define correspondingly  $S_0^{-i}$  and  $\Sigma_0^{-i}$ . Define for all  $n \geq 1$ ,  $\Sigma_n^i \equiv \{\sigma^i \in \Sigma^i : \sigma^i(s^i) > 0 \Rightarrow s^i \in S_{n-1}^i\}$  and  $S_n^i = S_{n-1}^i \setminus \{s^i \in S_{n-1}^i : \exists \sigma^i \in \Sigma_n^i \ni u_i(\sigma^i, \sigma^{-i}) > u_i(s^i, \sigma^{-i}), \forall \sigma^{-i} \in \Sigma_n^{-i}\}$ . Define  $S_\infty^i \equiv \bigcap_{n=1}^\infty S_n^i$ . Because  $S^A$  and  $S^B$  are finite sets, it takes a finite number of steps to reach  $S_\infty^i$ , for  $i = A, B$ . We must show that  $S_\infty^i$  is independent of the order of deletion of strictly dominated strategies.

dominated strictly by  $\sigma'_A$  and hence  $a_j$  cannot be contained in  $V_N^A$ , as long as  $\sigma'_A$  is feasible when A is restricted to using pure strategies in  $V_N^A$ . Note that if  $\sigma'_A$  is again infeasible when A is restricted to using pure strategies in  $V_N^A$ , then by repeating the above argument we can again create some  $\sigma''_A$  that dominates strictly  $\sigma'_A$ . In the end, since  $S^A$  is a finite set, there must be some  $\sigma_A^*$  that remains feasible when A is restricted to using pure strategies in  $V_N^A$ ,<sup>2</sup> and  $\sigma_A^*$  dominates strictly  $a_j$ , implying that  $V_N^A$  cannot contain  $a_j$ , which is a contradiction.

Thus we conclude that there must exist some  $b_k \in V_N^B \cap [T_m^B]^c$  such that when player B uses  $b_k$ , player A feels that using  $a_j$  is no worse than using  $\sigma_A$ . As  $T_m^B$  does not contain  $b_k$ , there exists  $\sigma_B$  that dominates strictly  $b_k$  and makes  $b_k$  deleted at some step  $l \leq m - 1$  under  $T$ . If  $V_N^A \subset T_l^A$ , then by mimicking the reasoning in the preceding paragraph, we can establish that  $V_N^B$  does not contain  $b_k$ , which would be a contradiction. Thus there must exist some  $a_h \in V_N^A \cap [T_l^A]^c$ , such that when player A uses  $a_h$ , player B feels that using  $b_k$  is no worse than using  $\sigma_B$ .

Note that  $a_j$  and  $a_h$  are distinct:  $a_j$  has not been deleted at step  $l \leq m - 1$  under  $T$ , and if  $a_j = a_h$ , then  $b_k$  cannot be deleted at step  $l$  under  $T$ !

Now we can turn the spotlight away from  $a_j \in V_N^A \cap [T_M^A]^c$  but on to  $a_h \in V_N^A \cap [T_l^A]^c$  instead. Define  $a_1 \equiv a_j$  and  $a_2 \equiv a_h$ , and repeat the above argument. We thus obtain an infinite sequence of distinct elements of  $S^A$ , which, given that  $S^A$  is finite, is a contradiction. We conclude that  $V_N^A \subset T_M^A$ , and by symmetry, that  $T_M^A \subset V_N^A$ , so that  $V_N^A = T_M^A$ . This conclusion applies to player B as well.

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<sup>2</sup>First note the following fact: If  $\sigma_A$  strictly dominates  $a_j$  when B is restricted to using pure strategies in  $S_B \subset S^B$ , then for all  $b_r \in S_B$ ,  $\sum_{i=1}^m u_A(a_i, b_r) \sigma_A(a_i) - u_A(a_j, b_r) > 0$ , implying that  $\sigma_A^j \equiv \frac{1}{1 - \sigma_A(a_j)} (\sigma_A(a_1), \sigma_A(a_2), \dots, \sigma_A(a_{j-1}), \sigma_A(a_{j+1}), \dots, \sigma_A(a_m))$  is also a mixed strategy that dominates strictly  $a_j$ . Note that  $\sigma_A^j$  assigns zero probability to  $a_j$ . Thus when  $a_j$  is strictly dominated, we can always assume that it is dominated by a mixed strategy that assigns positive probabilities only to A's pure strategies that differ from  $a_j$ . Now, define  $a^1$  as  $a_i$  mentioned above, and if  $\sigma'_A$  is infeasible when A is restricted to using pure strategies in  $V_N^A$ , then there is  $a^2 \in [V_N^A]^c$  with  $\sigma'_A(a^2) > 0$  and  $a^2$  differs from  $a^1$ . Repeating this argument, we see that if the afore-mentioned  $\sigma_A^*$  does not exist, then we would obtain an infinite sequence  $\{a^n\}$  contained in  $S^A \cap [V_N^A]^c$ , which contradicts the fact that  $S^A$  is finite.